Reasoning in Games

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Plan

- ✓ Day 1: Decision Theory
- ✓ Day 2: From Decisions to Games
 - Day 3: Game Models
 - Day 4: Modeling Deliberation (in Games)
 - Day 5: Backward and Forward Induction, Concluding Remarks (Language-Based Games/ Variable Frame Theory, Behavioral Game Theory, ...)



Strategic Game













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- A model of a game is a completion of the partial specification of the Bayesian decision problems and a representation of a particular play of the game.
- There are no special rules of rationality telling one what to do in the absence of degrees of belief except: decide what you believe, and then maximize (subjective) expected utility.

Suppose that G is a game.

- Outcomes of the game: $S = \prod_{i \in N} S_i$
- A profile is a vector $\vec{s} \in S$, specifying an action for each player
- ▶ Player *i*'s partial beliefs (or conjecture): $P_i \in \Delta(S_{-i})$

 $\Delta(X)$ is the set of probabilities measures over X

- ▶ W is a set of *possible worlds* (possible outcomes of the game)
- ► **s** is a function $\mathbf{s} : W \to \prod_{i \in N} S_i$, write $\mathbf{s}_i(w)$ for the *i*th component of $\mathbf{s}(w)$

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- ► **s** is a function **s** : $W \to \prod_{i \in N} S_i$, write $\mathbf{s}_i(w)$ for the *i*th component of $\mathbf{s}(w)$
- ▶ If $\vec{s} \in \prod_{i \in N} S_i$, then $[\vec{s}] = \{w \mid \mathbf{s}(w) = \vec{s}\}$; if $s_i \in S_i$, then $[s_i] = \{w \mid \mathbf{s}_i(w) = s_i\}$; and if $X \subseteq S$, $[X] = \bigcup_{s \in X} [s]$.

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- ex ante beliefs: For each i ∈ N, let P_i ∈ Δ(W) (the set of probability measures on W). Two assumptions:
 - [s] is measurable for all strategy profiles $s \in S$
 - $P_i([s_i]) > 0$ for all $s_i \in S_i$

ex interim beliefs: $P_{i,w} \in \Delta(S_{-i})$

- ...given player *i*'s choice: $P_{i,w}(\cdot) = P_i(\cdot | [\mathbf{s}_i(w)])$
- ▶ ...given all player *i* knows: $P_{i,w}(\cdot) = P_i(\cdot | K_i), \qquad K_i \subseteq [\mathbf{s}_i(w)]$
- ▶ ...given all player *i* fully believes: $P_{i,w}(\cdot) = P_i(\cdot | B_i), B_i \subseteq [\mathbf{s}_i(w)]$

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Expected utility of strategy $s_i \in S_i$: Given $P \in \Delta(S_{-i})$,

$$EU_{i,P}(s_i) = \sum_{s_{-i} \in S_{-i}} P(s_{-i})u_i(s_i, s_{-i})$$

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Expected utility of strategy $s_i \in S_i$: Given $w \in W$,

$$EU_{i,\mathbf{w}}(s_i) = \sum_{s_{-i} \in S_{-i}} P_{i,\mathbf{w}}([s_{-i}])u_i(s_i, s_{-i})$$























 Ann's choice is optimal (given her information)









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$$1 \cdot \frac{3}{4} + 0 \cdot \frac{1}{4} \ge 0 \cdot \frac{3}{4} + 2 \cdot \frac{1}{4}$$





- Ann's choice is optimal (given her information)
- Bob's choice is optimal (given her information)
- Bob considers it possible Ann is irrational

$$1 \cdot \frac{1}{2} + \mathbf{0} \cdot \frac{1}{2} \not\ge \mathbf{0} \cdot \frac{1}{2} + \mathbf{2} \cdot \frac{1}{2}$$

For any $P \in \Delta(S_{-i})$ and $s_i \in S_i$, $EU_{i,P}(s_i) = \sum_{s_{-i} \in S_{-i}} P(s_{-i})u_i(s_i, s_{-i})$
For any $w \in W$ and $s_i \in S_i$, $EU_{i,w}(s_i) = \sum_{s_{-i} \in S_{-i}} P_{i,w}([s_{-i}])u_i(s_i, s_{-i})$

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 $Rat_i = \{w \mid EU_{i,w}(\mathbf{s}_i(w)) \ge EU_{i,w}(\mathbf{s}_i) \text{ for all } \mathbf{s}_i \in S_i\}$

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Each $P \in \Delta(W)$ is associated with $P^S \in \Delta(S)$ as follows: for all $s \in S$, $P^S(s) = P([s])$

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A mixed strategy $\sigma \in \prod_{i \in N} \Delta(S_i)$, $P_{\sigma} \in \Delta(S)$, $P_{\sigma}(s) = \sigma_1(s_1) \cdots \sigma_n(s_n)$

Characterizing Nash Equilibria

Theorem (Aumann). σ is a Nash equilibrium of *G* iff there exists a model $\mathcal{M}^G = \langle W, (P_i)_{i \in N}, \mathbf{s} \rangle$ such that:

- for all $i \in N$, Rat_i = W;
- for all $i, j \in N$, $P_i = P_j$; and
- for all $i \in N$, $P_i^S = P_{\sigma}$.



Correlation: Players can improve their expected value by correlating their choices on an "outside signal"

Correlated Strategies

	L	R
U	2, 1	0, 0
D	0, 0	1, 2

- Three Nash equilibria:
 - (*U*, *L*): the payoff is (2, 1)
 - (*D*, *R*): the payoff is (1, 2)
 - $([\frac{2}{3}(U), \frac{1}{3}D], [\frac{1}{3}(L), \frac{2}{3}(R)])$: the payoff is $(\frac{2}{3}, \frac{2}{3})$

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Correlated Strategies

	L	R		L	R
U	2, 1	0, 0	U	0.5	0
D	0, 0	1, 2	D	0	0.5

- Three Nash equilibria:
 - (*U*, *L*): the payoff is (2, 1)
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 - $([\frac{2}{3}(U), \frac{1}{3}D], [\frac{1}{3}(L), \frac{2}{3}(R)])$: the payoff is $(\frac{2}{3}, \frac{2}{3})$
- Mixed Strategies: Each player conducts a private, independent lottery to choose their strategy.
- ► Conduct a *public* lottery: flip a fair coin and follow the strategy $(H \Rightarrow (U, L), T \Rightarrow (D, R))$. The payoff is (1.5, 1.5).

Two extremes:

- 1. Completely private, independent lotteries
- 2. A single, completely public lottery

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What about: a public lottery, but reveal only partial information about the outcome to each of the players?

	L	R		
U	6, 6	2, 7		
D	7, 2	0, 0		

- Three Nash equilibria:
 - (*U*, *R*): the payoff is (2, 7)

 - (D, L): the payoff is (7,2) $([\frac{2}{3}(U), \frac{1}{3}D], [\frac{2}{3}(L), \frac{1}{3}(R)])$: the payoff is $(4\frac{2}{3}, 4\frac{2}{3})$

	L	R		L	R
U	6, 6	2, 7	U	1/3	1/3
D	7, 2	0, 0	D	1/3	0

- Three Nash equilibria:
 - (*U*, *R*): the payoff is (2, 7)
 - (*D*, *L*): the payoff is (7, 2)
 - $([\frac{2}{3}(U), \frac{1}{3}D], [\frac{2}{3}(L), \frac{1}{3}(R)])$: the payoff is $(4\frac{2}{3}, 4\frac{2}{3})$
- After conducting the lottery, an outside observer provides Ann with a recommendation to play the first component of the profile that was chosen, and Bob the second component.

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- After conducting the lottery, an outside observer provides Ann with a recommendation to play the first component of the profile that was chosen, and Bob the second component.
- ► The expected payoff is ¹/₃(6,6) + ¹/₃(2,7) + ¹/₃(7,2) = (5,5) (which is outside the convex hull of the Nash equilibria)



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With more than 2 players...

- A player may believe that (some of) the other players strategy choices are **independent** or **correlated**.
- Two players can agree or disagree on the probabilities that the assign to a third player's choice of strategy.

Characterizing Correlated Equilibrium

Theorem (Aumann). σ is a correlated equilibrium of *G* iff there exists a model $\mathcal{M}^G = \langle W, (P_i)_{i \in N}, \mathbf{s} \rangle$ such that:

- for all $i \in N$, Rat_i = W; and
- for all $i \in N$, $P_i^S = \sigma$.

Rationalizability

A **best reply set** (BRS) is a sequence $(B_1, B_2, ..., B_n) \subseteq S = \prod_{i \in N} S_i$ such that for all $i \in N$, and all $s_i \in B_i$, there exists $\mu_{-i} \in \Delta(B_{-i})$ such that s_i is a best response to μ_{-i} : I.e.,

$$b_i = rg\max_{s_i \in S_i} EU_{i,\mu_{-i}}(s_i)$$

		b ₁	b ₂	b ₃	b_4
	a ₁	0, 7	2, 5	7, 0	0, 1
1	a 2	5, 2	3, 3	5, 2	0, 1
	a 3	7, 0	2, 5	0, 7	0, 1
	<i>a</i> ₄	0, 0	0, -2	0, 0	10, -1

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	2									
		<i>b</i> ₁	b ₂	b ₃	b ₄					
	a ₁	0, 7	2, 5	7, 0	0, 1					
1	a 2	5, 2	3, 3	5, 2	0, 1					
	a 3	7, 0	2, 5	0, 7	0, 1					
	<i>a</i> ₄	0, 0	0, -2	0, 0	10, -1					

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- ► (a_2, b_2) is the unique Nash equilibria, hence $(\{a_2\}, \{b_2\})$ is a BRS
- $(\{a_1, a_3\}, \{b_1, b_3\})$ is a BRS
- $(\{a_1, a_2, a_3\}, \{b_1, b_2, b_3\})$ is a full BRS

Theorem (Bernheim; Pearce; Brandenburger and Dekel; ...). $(B_1, B_2, ..., B_n)$ is a BRS for *G* iff there exists a model $\mathcal{M}^G = \langle W, (P_i)_{i \in N}, \mathbf{s} \rangle$ such that for all $i \in N$, $\operatorname{Rat}_i = W$ and $[B_1 \times \cdots \times B_n] = W$.







Game 2: U strictly dominates D



Game 2: *U* strictly dominates *D*, and *after removing D*, *L* strictly dominates *R*.



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Theorem. In all models where the players are *rational* and there is *common belief of rationality*, the players choose strategies that survive iterative removal of strictly dominated strategies (and, conversely...).

Comparing Dominance Reasoning and MEU

 $G = \langle N, (S_i)_{i \in N}, (u_i)_{i \in N} \rangle$

 $X \subseteq S_{-i}$ (a set of strategy profiles for all players except *i*)

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 $s, s' \in S_i$, s strictly dominates s' with respect to X provided

$$\forall s_{-i} \in X, \ u_i(s, s_{-i}) > u_i(s', s_{-i})$$

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 $s, s' \in S_i$, s strictly dominates s' with respect to X provided

$$\forall s_{-i} \in X, \ u_i(s, s_{-i}) > u_i(s', s_{-i})$$

 $p \in \Delta(X)$, *s* is a **best response** to *p* with respect to *X* provided $\forall s' \in S_i, EU(s, p) \ge EU(s', p)$



D is strictly dominated by (0.5U, 0.5M).

Strict Dominance and MEU

Proposition. Suppose that $G = \langle N, (S_i)_{i \in N}, (u_i)_{i \in N} \rangle$ is a strategic game and $X \subseteq S_{-i}$. A strategy $s_i \in S_i$ is strictly dominated (possibly by a mixed strategy) with respect to X iff there is no probability measure $p \in \Delta(X)$ such that s_i is a best response to p.

Important Issue: Correlated Beliefs



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Note that y is not strictly dominated for Charles.

Important Issue: Correlated Beliefs

X		r	y	Ι	r	Ζ	Ι	r
и	1,1,3	1,0,3	и	1,1,2	1,0,0	и	1,1,0	1,0,0
d	0,1,0	0,0,0	d	0,1,0	1,1,2	d	0,1,3	0,0,3

- Note that y is not strictly dominated for Charles.
- ▶ It is easy to find a probability measure $p \in \Delta(S_A \times S_B)$ such that *y* is a best response to *p*. Suppose that $p(u, l) = p(d, r) = \frac{1}{2}$. Then, EU(x, p) = EU(z, p) = 1.5 while EU(y, p) = 2.
Important Issue: Correlated Beliefs

X		r	y	Ι	r	Ζ		r
и	1,1,3	1,0,3	и	1,1,2	1,0,0	и	1,1,0	1,0,0
d	0,1,0	0,0,0	d	0,1,0	1,1,2	d	0,1,3	0,0,3

- Note that y is not strictly dominated for Charles.
- ▶ It is easy to find a probability measure $p \in \Delta(S_A \times S_B)$ such that *y* is a best response to *p*. Suppose that $p(u, l) = p(d, r) = \frac{1}{2}$. Then, EU(x, p) = EU(z, p) = 1.5 while EU(y, p) = 2.
- However, there is no probability measure p ∈ Δ(S_A × S_B) such that y is a best response to p and p(u, l) = p(u) · p(l).

X		r	у	Ι	r	Ζ	Ι	r
u	1,1,3	1,0,3	u	1,1,2	1,0,0	 и	1,1,0	1,0,0
d	0,1,0	0,0,0	d	0,1,0	1,1,2	d	0,1,3	0,0,3

To see this, suppose that a is the probability assigned to u and b is the probability assigned to I. Then, we have:

• The expected utility of y is 2ab + 2(1-a)(1-b);

- The expected utility of *x* is 3*ab* + 3*a*(1 − *b*) = 3*a*(*b* + (1 − *b*)) = 3*a*; and
- The expected utility of z is

$$3(1-a)b + 3(1-a)(1-b) = 3(1-a)(b+(1-b)) = 3(1-a).$$

Let $P \in \Delta(X)$ be a probability measure, the **support** of *P* is $supp(P) = \{x \in X \mid P(x) > 0\}.$

A probability measure $P \in \Delta(X)$ is said to be a **full support** probability measure on *X* provided supp(P) = X.



Is R rationalizable?



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There is no full support probability such that R is a best response



Is R rationalizable?

There is no *full support* probability such that R is a best response Should Ann assign probability 0 to R or probability > 0 to R?

Strategic Reasoning and Admissibility

"The argument for deletion of a weakly dominated strategy for player *i* is that he contemplates the possibility that every strategy combination of his rivals occurs with positive probability. However, this hypothesis clashes with the logic of iterated deletion, which assumes, precisely, that eliminated strategies are not expected to occur."

Mas-Colell, Whinston and Green. Introduction to Microeconomics. 1995.

A Puzzle

R. Cubitt and R. Sugden. *Rationally Justifiable Play and the Theory of Non-cooperative games.* Economic Journal, 104, pgs. 798 - 803, 1994.

R. Cubitt and R. Sugden. *Common reasoning in games: A Lewisian analysis of common knowledge of rationality.* Manuscript, 2011.







Game 2: U weakly dominates D



Game 2: *U* weakly dominates *D*, and *after removing D*, *L* strictly dominates *R*.



Game 2: But, now what is the reason for not playing *D*?



Game 2: But, now what is the reason for not playing *D*?

Theorem (Samuelson). There is no model of Game 2 satisfying common knowledge of rationality (where rationality incorporates weak dominance).



There is no model of this game with common knowledge of admissibility.



The "full" model of the game



The "full" model of the game: B is not admissible given Ann's information



What is wrong with this model?



Privacy of Tie-Breaking/No Extraneous Beliefs: If a strategy is *rational* for an opponent, then it cannot be "ruled out".



► Game models describe the *informational context* of a game.

 Game models can be used to *characterize* different solution concepts (e.g., iterated strict dominance, iterated weak dominance, Nash equilibrium, correlated equilibrium,...)

Tomorrow: Deliberation in Games