

Probabilistic Methods in Social Choice

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Voters Rankings

1

a b c d

2

b a d c

3

b d a c

4

d c a b

Aggregation Method



Winner

Winning set

Defeat Relation

Committee

Voters Rankings

1 *a b c d*

2 *b a d c*

3 *b d a c*

4 *d c a b*

Axiomatic
Characterization

Winner

Winning set

Defeat Relation

Committee

Voters Rankings

1	<i>a b c d</i>
2	<i>b a d c</i>
3	<i>b d a c</i>
4	<i>d c a b</i>

Aggregation Method

Winner
Winning set
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1. Probabilistically generate preferences

2. Probability judgements

Voters Rankings

1 *a b c d*

2 *b a d c*

3 *b d a c*

4 *d c a b*

Aggregation Method

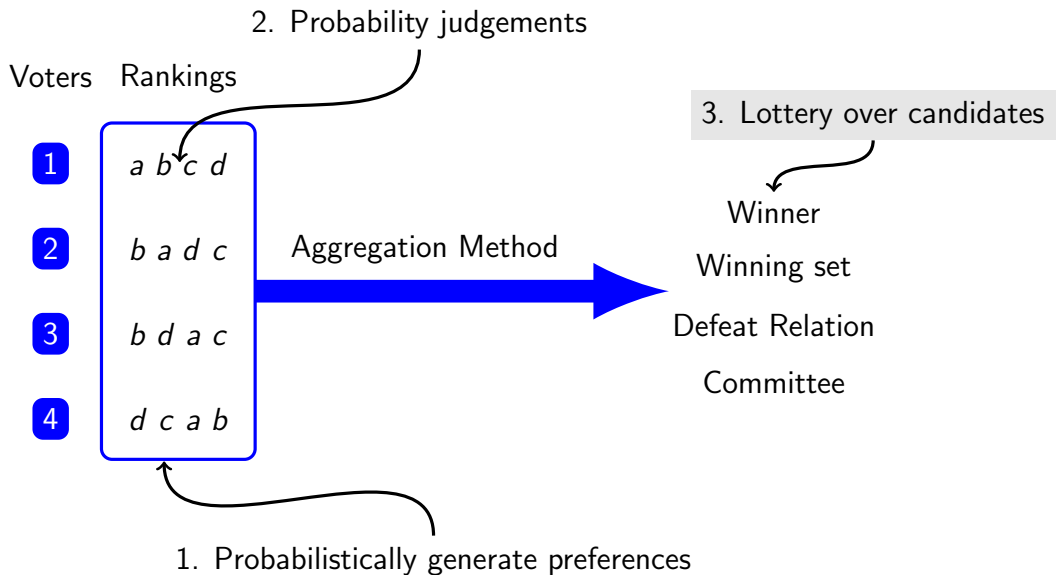
Winner

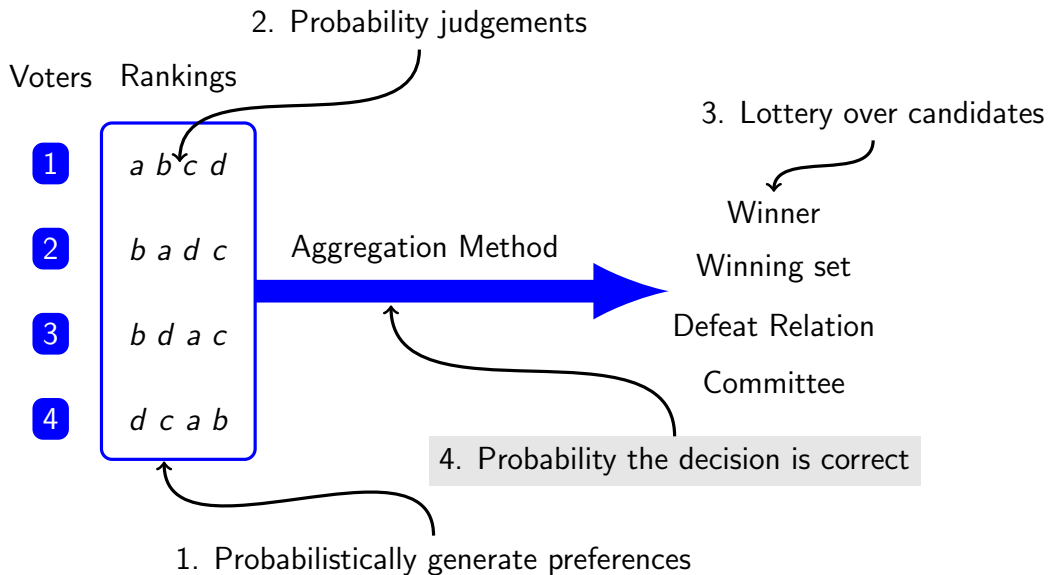
Winning set

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1. Probabilistically generate preferences





Plan

- ▶ Background on voting theory
- ▶ Generating preference profiles
- ▶ Quantitative analysis of voting methods
- ▶ Probabilistic voting methods
- ▶ Condorcet jury theorem and related results
- ▶ Aggregating probabilistic judgements

Voting

New York City Voters Just Adopted Ranked-Choice Voting in Elections. Here's How It Works



Kyle Winkler/Reuters for the Wall Street Journal via Getty Images/Getty Images for New York in Blue. © 2020

The Rules of the Game: A New Electoral System

By Michael and Joseph Katz

January 16, 2017 10:42



Opinion

Article

How Majority Rule Might Have Stopped Donald Trump

By Michael and Joseph Katz

January 16, 2017

5 4 3 2 1

100 90 80 70 60 50 40 30 20 10 0

Donald Trump's victory in the 2016 election was a surprise to many. But it was not a surprise to those who have studied the history of the electoral college.

There is a reason that the electoral college has been so controversial. It is a system that was designed to protect the interests of the small states. But it has also been a source of controversy because it has allowed a candidate to win the presidency without winning the popular vote.

One of the most interesting aspects of the electoral college is that it has allowed a candidate to win the presidency without winning the popular vote. This has happened in five of the 57 presidential elections in U.S. history.



<https://www.electology.org>

<http://www.fairvote.org>

Rankings

Let X be a set of candidates and V a set of voters.

A **ranking** of X is a strict linear order P on X : a relation $P \subseteq X \times X$ satisfying the following conditions for all $x, y, z \in X$:

asymmetry: if $x P y$ then *not* $y P x$;

transitivity: if $x P y$ and $y P z$, then $x P z$;

weak completeness: if $x \neq y$, then $x P y$ or $y P x$.

Let $\mathcal{L}(X)$ be the set of all strict linear orders on X .

Rankings

SANTA CLARA COUNTY						
MAYOR 市長	1 1st Choice 第一選擇	2 2nd Choice 第二選擇	3 3rd Choice 第三選擇	4 4th Choice 第四選擇	5 5th Choice 第五選擇	6 6th Choice 第六選擇
ELLEN LEE ZHOU / 李麗晨 Behavioral Health Clinician 行為健康臨床治療師	● 1	2	3	4	5	6
LONDON N. BREED / 倫敦·布里德 Mayor of San Francisco 三藩市市長	1	2	● 3	4	5	6
JOEL VENTRESCA / 喬爾·范崔斯卡 Retired Airport Analyst 退休機場分析師	1	2	3	4	● 5	6
WILMA PANG / 彭德慧 Retired Music Professor 退休音樂教授	1	2	3	● 4	5	6
ROBERT L. JORDAN, JR. / 小羅伯特·L·喬丹 Preacher 傳教士	1	2	3	4	5	● 6
PAUL YBARRA ROBERTSON / 保羅·伊巴拉·羅伯森 Small Business Owner 小企業業主	1	● 2	3	4	5	6
	1	2	3	4	5	6

Variable candidate/voter profiles

Fix infinite sets \mathcal{V} and \mathcal{X} of *voters* and *candidates*, respectively.

For $X \subseteq \mathcal{X}$, let $\mathcal{L}(X)$ be the set of all strict linear orders on X .

A *profile* is a function $\mathbf{P} : V(\mathbf{P}) \rightarrow \mathcal{L}(X(\mathbf{P}))$ for some nonempty finite $V(\mathbf{P}) \subseteq \mathcal{V}$ and nonempty finite $X(\mathbf{P}) \subseteq \mathcal{X}$.

We call $V(\mathbf{P})$ and $X(\mathbf{P})$ the sets of *voters in \mathbf{P}* and *candidates in \mathbf{P}* , respectively.

We call $\mathbf{P}(i)$ voter i 's *ranking*, and we write ' $x\mathbf{P}_iy$ ' for $(x, y) \in \mathbf{P}(i)$. As usual, we take $x\mathbf{P}_iy$ to mean that voter i strictly prefers candidate x to candidate y .

Anonymous profiles

40	35	25
<i>t</i>	<i>r</i>	<i>k</i>
<i>k</i>	<i>k</i>	<i>r</i>
<i>r</i>	<i>t</i>	<i>t</i>

Margin

Let \mathbf{P} be a profile and $a, b \in X(\mathbf{P})$. Then the **margin of a over b** is:

$$\text{Margin}_{\mathbf{P}}(a, b) = |\{i \in V(\mathbf{P}) \mid a \mathbf{P}_i b\}| - |\{i \in V(\mathbf{P}) \mid b \mathbf{P}_i a\}|.$$

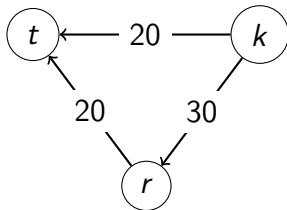
We say that a is **majority preferred** to b in \mathbf{P} when $\text{Margin}_{\mathbf{P}}(a, b) > 0$.

Margin Graph

The **margin graph of P** , $\mathcal{M}(P)$, is the weighted directed graph whose set of nodes is $X(P)$ with an edge from a to b weighted by $\text{Margin}(a, b)$ when $\text{Margin}(a, b) > 0$. We write

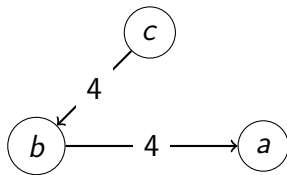
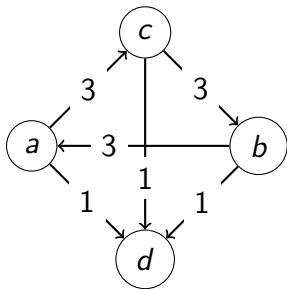
$$a \xrightarrow{\alpha}_P b \text{ if } \alpha = \text{Margin}_P(a, b) > 0.$$

40	35	25
t	r	k
k	k	r
r	t	t



Margin Graph

A **margin graph** is a weighted directed graph \mathcal{M} where all the weights have the same parity.



Theorem (Debord, 1987)

For any margin graph \mathcal{M} , there is a profile \mathbf{P} such that \mathcal{M} is the margin graph of \mathbf{P} .

A **voting method** is a function F on the domain of all profiles such that for any profile P , $\emptyset \neq F(P) \subseteq X(P)$ (also called a **variable social choice correspondence** VSCC).

A **variable-election collective choice rule** (VCCR) is a function f on the domain of all profiles such that for any profile P , $f(P)$ is an asymmetric binary relation on $X(P)$, which we call *the defeat relation for P under f* .

For $x, y \in X(P)$, we say that x *defeats y in P according to f* when $(x, y) \in f(P)$.

Majority Defeat

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Example

In the 2000 U.S. presidential election in Florida, George W. Bush defeated Al Gore and Ralph Nader according to Plurality voting, which only allows voters to vote for one candidate. Yet assuming that most Nader voters preferred Gore to Bush, it follows that a majority of all voters preferred Gore to Bush.

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Under Plurality, Nader spoiled the election for Gore, handing victory to Bush.

Anonymity and Neutrality

Anonymity: if x defeats y in P , and P' is obtained from P by swapping the ballots assigned to two voters, then x still defeats y in P' .

Neutrality: if x defeats y in P , and P' is obtained from P by swapping x and y on each voter's ballot, then y defeats x in P' .

Availability: for all profiles P , there is some undefeated candidate.

Monotonicity

Monotonicity (resp. Monotonicity for two-candidate profiles): if x defeats y in a profile (resp. two-candidate profile) P , and P' is obtained from P by some voter i moving x above the candidate that i ranked immediately above x in P , then x defeats y in P' .

Lemma

If f satisfies Anonymity, Neutrality, and Monotonicity with respect to two-candidate profiles, then f satisfies Special Majority Defeat: for any two-candidate profile P , x defeats y in P according to f only if x is majority preferred to y .

Other rules satisfying Anonymity, Neutrality and Monotonicity: The completely indecisive method; Unanimity; Quota rules (cf. Fishburn 1974, Section 1)

Neutral Reversal: if P' is obtained from P by adding two voters with reversed ballots, then x defeats y in P if and only if x defeats y in P' .

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Pareto: if for all profiles P and $x, y \in X(P)$, if $xP_i y$ for all $i \in V(P)$, then x defeats y in P .

Characterizing Majority Rule

Proposition

For any VCCR f , the following are equivalent:

- 1. f coincides with majority rule on two-candidate profiles;*
- 2. f satisfies the following axioms with respect to two-candidate profiles:
Anonymity, Neutrality, Monotonicity, Pareto, and Upward Neutral Reversal.*

W. Holliday and EP. *Axioms for Defeat in Democratic Elections*. Forthcoming Journal of Theoretical Politics, <https://arxiv.org/pdf/2008.08451.pdf>, 2021.

Characterizing Majority Rule

K. May. *A Set of Independent Necessary and Sufficient Conditions for Simple Majority Decision*. *Econometrica*, Vol. 20 (1952).

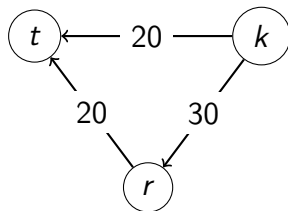
G. Asan and R. Sanver. *Another Characterization of the Majority Rule*. *Economics Letters*, 75 (3), 409-413, 2002.

E. Maskin. *Majority rule, social welfare functions and game forms*. in *Choice, Welfare and Development*, The Clarendon Press, pgs. 100 - 109, 1995.

G. Woeginger. *A new characterization of the majority rule*. *Economic Letters*, 81, pgs. 89 - 94, 2003.

More than 2 candidates...

40	35	25
<i>t</i>	<i>r</i>	<i>k</i>
<i>k</i>	<i>k</i>	<i>r</i>
<i>r</i>	<i>t</i>	<i>t</i>



Plurality winner	<i>t</i>
Instant Runoff winner	<i>r</i>
Borda winner	<i>k</i>
Condorcet winner	<i>k</i>
Condorcet loser	<i>t</i>

Positional scoring rules

A *scoring vector* is a vector $\langle s_1, \dots, s_n \rangle$ of numbers such that for each $m \in \{1, \dots, n-1\}$, $s_m \geq s_{m+1}$.

Given a profile \mathbf{P} with $|X(\mathbf{P})| = n$, $x \in X(\mathbf{P})$, a scoring vector \vec{s} of length n , and $i \in V(\mathbf{P})$, define $score_{\vec{s}}(x, \mathbf{P}_i) = s_r$ where $r = Rank(x, \mathbf{P}_i)$.

Let $score_{\vec{s}}(x, \mathbf{P}) = \sum_{i \in V(\mathbf{P})} score_{\vec{s}}(x, \mathbf{P}_i)$. A voting method F is a **positional scoring rule** if there is a map \mathcal{S} assigning to each natural number n a scoring vector of length n such that for any profile \mathbf{P} with $|X(\mathbf{P})| = n$,

$$F(\mathbf{P}) = \operatorname{argmax}_{x \in X(\mathbf{P})} score_{\mathcal{S}(n)}(x, \mathbf{P}).$$

Examples

Borda: $\mathcal{S}(n) = \langle n-1, n-2, \dots, 1, 0 \rangle$

Plurality: $\mathcal{S}(n) = \langle 1, 0, \dots, 0 \rangle$

Anti-Plurality: $\mathcal{S}(n) = \langle 1, 1, \dots, 1, 0 \rangle$

1	3	2	4
<i>a</i>	<i>b</i>	<i>b</i>	<i>c</i>
<i>c</i>	<i>a</i>	<i>c</i>	<i>a</i>
<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>

Borda winner *c*

Plurality winner *b*

Anti-Plurality winner *a*

Plurality vs. Borda

1	1
<hr/>	
<i>a</i>	<i>c</i>
<i>b</i>	<i>b</i>
<i>c</i>	<i>a</i>

Plurality winners: *a, b*

Borda winners: *a, b, c*

Iterative procedures: Instant Runoff

- ▶ If some alternative is ranked first by an absolute majority of voters, then it is declared the winner.
- ▶ Otherwise, the alternative ranked first by the fewest voters (the plurality loser) is eliminated.
- ▶ Votes for eliminated alternatives get transferred: delete the removed alternatives from the ballots and “shift” the rankings (e.g., if 1st place alternative is removed, then your 2nd place alternative becomes 1st).

Also known as Ranked-Choice, STV, Hare

How should you deal with ties? (e.g., multiple alternatives are plurality losers)

Non-neutral tiebreaking: Fix a linear ordering of the candidates

Remove all: Remove all candidates tied for the smallest plurality score

Parallel universe tiebreaking: A candidate a wins if a wins according to some linear ordering of the candidates

1	3	2	1	1
<hr/>				
c	c	b	a	a
a	b	a	c	b
b	a	c	b	c

Instant Runoff: $\{c\}$

Instant Runoff PUT: $\{a, c\}$

S. Obraztsova, E. Elkind and N. Hazon. *Ties Matter: Complexity of Voting Manipulation Revisited*. Proceedings of the Twenty-Second International Joint Conference on Artificial Intelligence.

J. Wang, S. Sikdar, T. Shepherd, Z. Zhao, C. Jiang and L. Xia. *Practical Algorithms for Multi-Stage Voting Rules with Parallel Universes Tiebreaking*. Proceedings of AAAI, 2019.

Iterative procedures

Variants:

- ▶ Plurality with runoff: remove all candidates except top two plurality score;
- ▶ Coombs: remove candidates with most last place votes;
- ▶ Baldwin: remove candidate with smallest Borda score;
- ▶ Nanson: remove candidates with below average Borda score

Example

1	1	1	1	1
<hr/>				
<i>c</i>	<i>b</i>	<i>a</i>	<i>b</i>	<i>d</i>
<i>a</i>	<i>d</i>	<i>b</i>	<i>c</i>	<i>a</i>
<i>d</i>	<i>a</i>	<i>c</i>	<i>d</i>	<i>b</i>
<i>b</i>	<i>c</i>	<i>d</i>	<i>a</i>	<i>c</i>

Instant Runoff $\{b\}$

Plurality with Runoff $\{a, b\}$

Coombs $\{d\}$

Baldwin $\{a, b, d\}$

Strict Nanson $\{a\}$