

Logics of Action, Ability, Knowledge and Obligation

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pacuit.org/esslli2019/epstit

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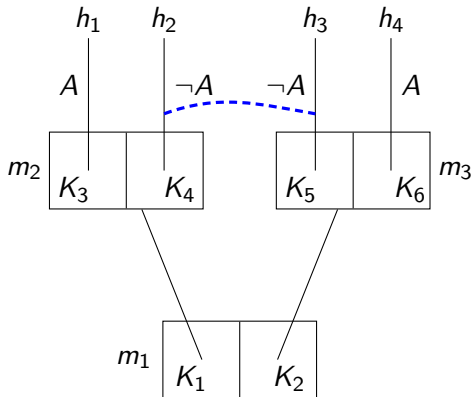
A. Herzig. *Logics of knowledge and action: critical analysis and challenges*. Autonomous Agent and Multi-Agent Systems, 2014.

V. Goranko and EP. *Temporal aspects of the dynamics of knowledge*. in Johan van Benthem on Logic and Information Dynamics, Outstanding Contributions to Logic, (eds. Alexandru Baltag and Sonja Smets), pp. 235 - 266, 2014.

J. Broeresen, A. Herzig and N. Troquard. *What groups do, can do and know they can do: An analysis in normal modal logics*. Journal of Applied and Non-Classical Logics, 19:3, pgs. 261 - 289, 2009.

W. van der Hoek and M. Wooldridge. *Cooperation, knowledge and time: Alternating-time temporal epistemic logic and its applications*. Studia Logica, 75, pgs. 125 - 157, 2003.

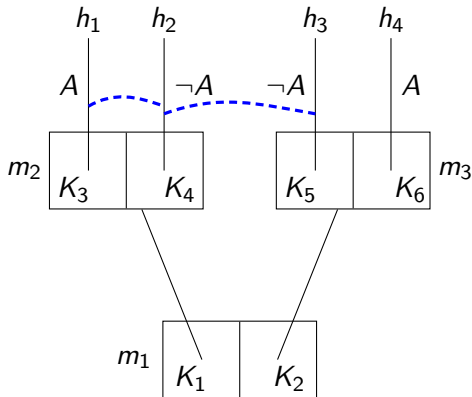
$\langle \text{Tree}, <, \text{Agent}, \text{Choice}, \{\sim_\alpha\}_{\alpha \in \text{Agent}}, V \rangle$



\sim_α is an equivalence relation on indices

$m/h \sim_\alpha m'/h'$: nothing α knows distinguishes m/h from m'/h' , or m/h and m'/h' are indistinguishable

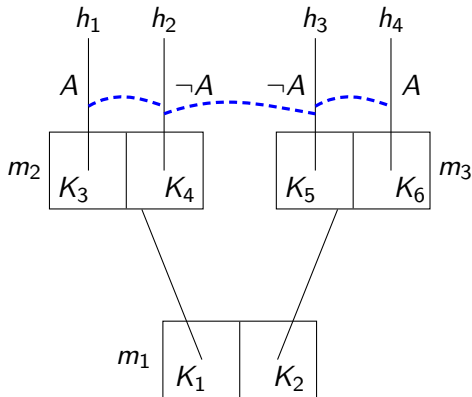
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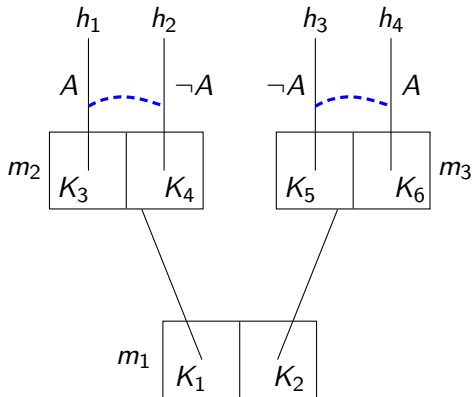
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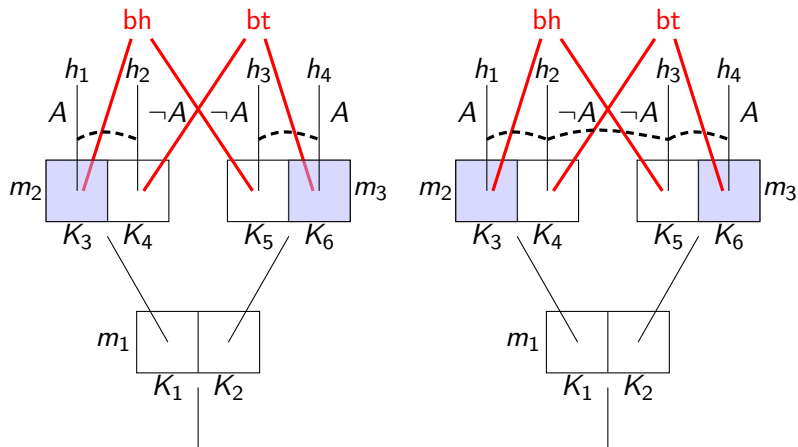
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Ability



Deliberative perspective

(C5) If $m/h \sim_{\alpha} m'/h'$, then $m/h'' \sim_{\alpha} m'/h'''$ for all $h'' \in H^m$ and $h''' \in H^{m'}$

Indistinguishability between moments: $m \sim_{\alpha} m'$ iff $m/h \sim_{\alpha} m'/h'$ for all $h \in H^m$ and $h' \in H^{m'}$.

Game Theory

A **game** is a mathematical model of a strategic interaction that includes

- ▶ the actions the players *can* take
- ▶ the players' interests (i.e., preferences),
- ▶ the “structure” of the decision problem

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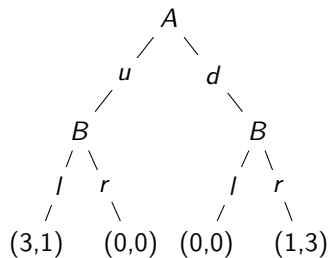
It does not specify the actions that the players do take.

Games

		<i>B</i>	
		l	r
<i>A</i>	u	3, 1	0, 0
	d	0, 0	1, 3

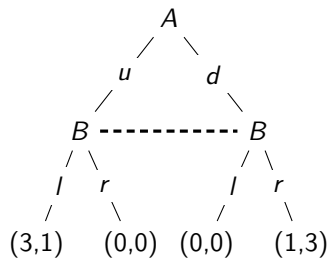
Games

		<i>B</i>	
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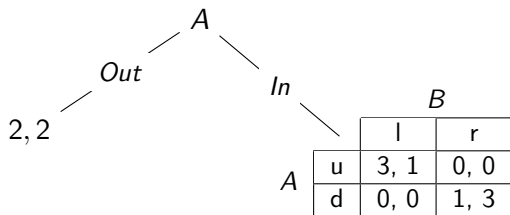
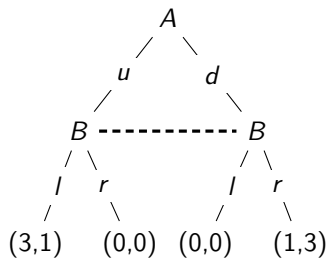
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		<i>B</i>	
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Knowledge and beliefs in game situations

J. Harsanyi. *Games with incomplete information played by "Bayesian" players I-III. Management Science Theory* **14**: 159-182, 1967-68.

R. Aumann. *Interactive Epistemology I & II. International Journal of Game Theory* (1999).

P. Battigalli and G. Bonanno. *Recent results on belief, knowledge and the epistemic foundations of game theory. Research in Economics* (1999).

R. Myerson. *Harsanyi's Games with Incomplete Information. Special 50th anniversary issue of Management Science*, 2004.

John C. Harsanyi, nobel prize winner in economics, developed a theory of games with **incomplete information**.

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John C. Harsanyi, nobel prize winner in economics, developed a theory of games with **incomplete information**.

1. **incomplete information**: uncertainty about the *structure* of the game (outcomes, payoffs, strategy space)
2. **imperfect information**: uncertainty *within the game* about the previous moves of the players

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- ▶ Varieties of informational attitudes
 - hard (“knowledge”)
 - soft (“beliefs”)

Ex ante vs. ex interim knowledge

- ▶ $\mathcal{M}, m/h \models K_\alpha A$ if and only if, for all m'/h' , if $m/h \sim_\alpha m'/h'$, then $\mathcal{M}, m'/h' \models A$
- ▶ $\mathcal{M}, m/h \models K_\alpha^{\text{act}} A$ if and only if, for all m'/h' , if $m/h \sim_\alpha m'/h'$ and $h' \in [\text{Type}_\alpha^m(h)]^{m'}$, $\mathcal{M}, m'/h' \models A$

Discussion

- ▶ Language/validities

$$\Box A \supset [\alpha \textit{ stit}: A]$$

$$K_\alpha \Box A \supset [\alpha \textit{ kstit}: A]$$

$$[\alpha \textit{ kstit}: A] \equiv K_\alpha^{\text{act}}[\alpha \textit{ stit}: A]$$

...

- ▶ What do the agents know vs. What do the agents know *given what they are doing*.

Broersen and Duijf

$m/h \sim'_\alpha m'/h'$ if and only if, $m/h \sim_\alpha m'/h'$ and $h' \in [Type_\alpha^m(h)]_\alpha^{m'}$

Theorem 1. *Let \mathcal{M} be a static labelled stit model, and let \mathcal{M}' be the associated transform epistemic stit model. Let $i \in \text{Ags}$ be an agent, $\varphi \in \mathcal{L}_{\text{stit}}$ be a standard stit formula, and m/h be an index. Then the following holds*

(1) $\mathcal{M}, m/h \models [i \text{ kstit}] \varphi$ if and only if $\mathcal{M}', m/h \models K_i [i \text{ stit}] \varphi$;

(2) $\mathcal{M}, m/h \models K_i \varphi$ if and only if $\mathcal{M}', m/h \models K_i \Box \varphi$;

and

(3) $\mathcal{M}', m/h \models K_i \varphi \rightarrow K_i [i \text{ stit}] \varphi$;

(4) $\mathcal{M}', m/h \models \Diamond K_i \varphi \rightarrow K_i \Diamond \varphi$

Concluding Remarks

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(cf. C. List, “Three kinds of collective attitudes”, 2015)
 - Group action types: products of individual action types vs. labelings of products of individual action tokens
 - Group oughts (Kooi and Tamminga: relativized to the interests of other groups of agents)

Conclude Remarks

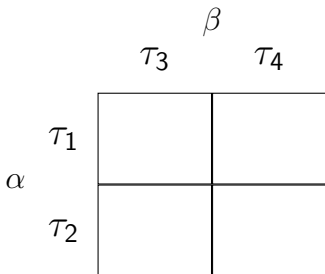
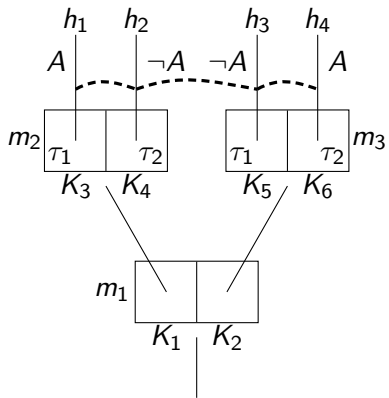
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- ▶ Moving beyond common payoffs
- ▶ When do two labeled stit models represent the *same* situation? (cf. when are two *games* the same?)



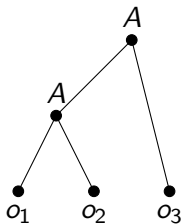
Thompson Transformations

Game-theoretic analysis should not depend on “irrelevant” features of the (mathematical) description of the game.

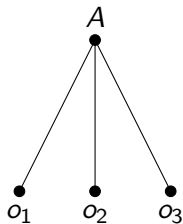
F. B. Thompson. *Equivalence of Games in Extensive Form*. Classics in Game Theory, pgs 36 - 45, 1952.

(Osborne and Rubinstein, pgs. 203 - 212)

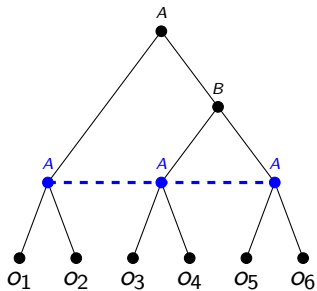
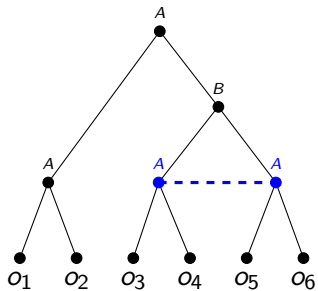
The same decision problem



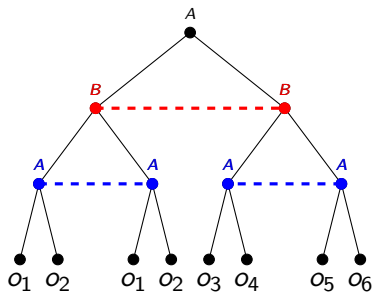
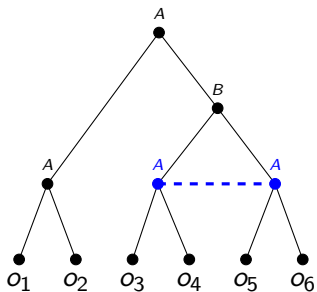
D_1



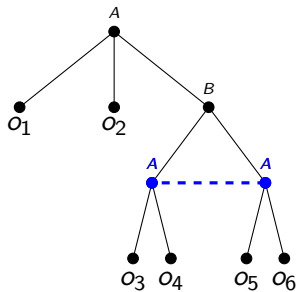
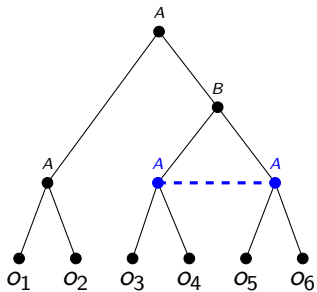
D_2



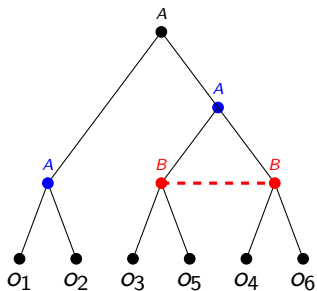
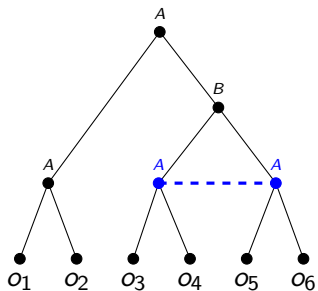
Inflation/deflation



Addition of Superfluous Move



Coalescing of moves



Interchange of moves

Theorem (Thompson) Each of the previous transformations preserves the reduced strategic form of the game. In finite extensive games (without uncertainty between subhistories), if any two games have the same reduced normal form then one can be obtained from the other by a sequence of the four transformations.

Extensive vs. Normal Forms

G. Bonanno. *Set-Theoretic Equivalence of Extensive-Form Games*. IJGT (1992).

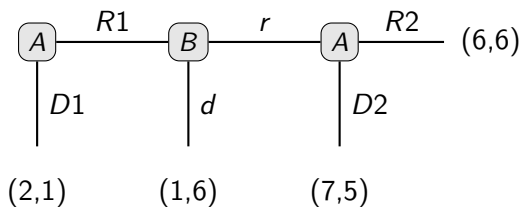
S. Elmes and P. Reny. *On The Strategic Equivalence of Extensive Form Games*. Journal of Economic Theory (1994).

E. Kholberg and F. Mertens. *On Strategic Stability of Equilibria*. Econometrica (1986).

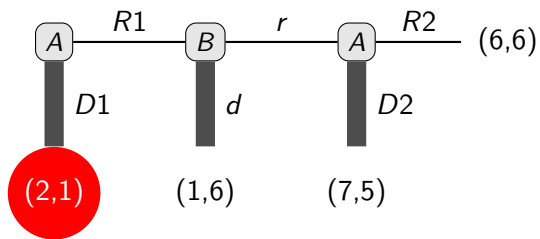
T. Seidenfeld. *When Normal and Extensive Form Decisions Differ*. Logic, Methodology and Philosophy of Science, 1994.

Strategic stit/obligations and incorporating game-theoretic reasoning (and beliefs)

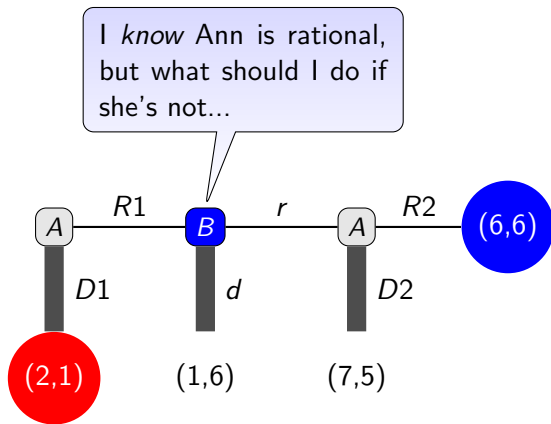
BI Puzzle?

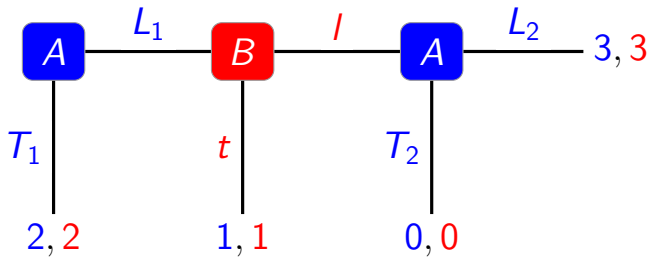


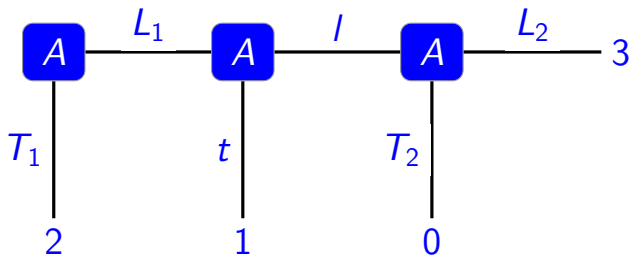
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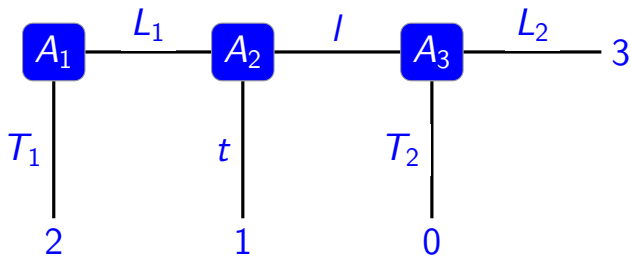


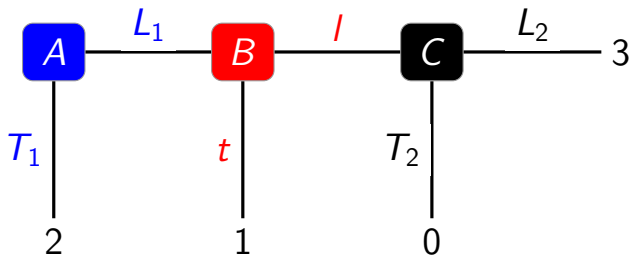
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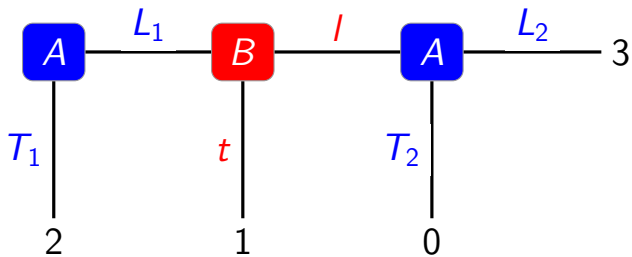












R. Aumann. *Backwards induction and common knowledge of rationality*. Games and Economic Behavior, 8, pgs. 6 - 19, 1995.

R. Stalnaker. *Knowledge, belief and counterfactual reasoning in games*. Economics and Philosophy, 12, pgs. 133 - 163, 1996.

J. Halpern. *Substantive Rationality and Backward Induction*. Games and Economic Behavior, 37, pp. 425-435, 1998.

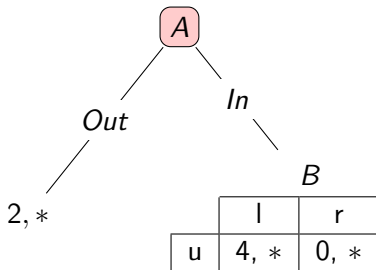
Informal characterizations of BI

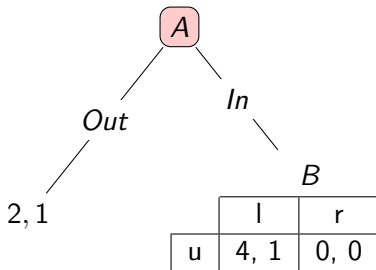
- ▶ Future choices are *epistemically independent* of any observed behavior
- ▶ Any “off-equilibrium” choice is interpreted simply as a mistake (which will not be repeated)
- ▶ At each choice point in a game, the players only reason about future paths

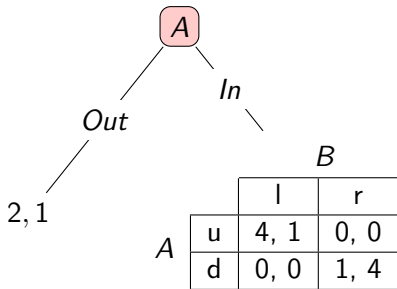
Rationalizing Observed Actions

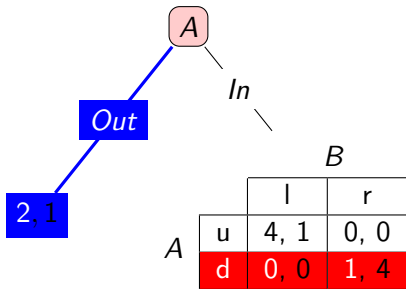
After observing an (unexpected) move by some player, you could:

1. Change your belief about the player's rationality, but maintain your beliefs about the player's *passive beliefs*.
2. Change your belief about the player's passive beliefs, but maintain your belief in the player's rationality.
3. Conclude that the player perceives the game differently.









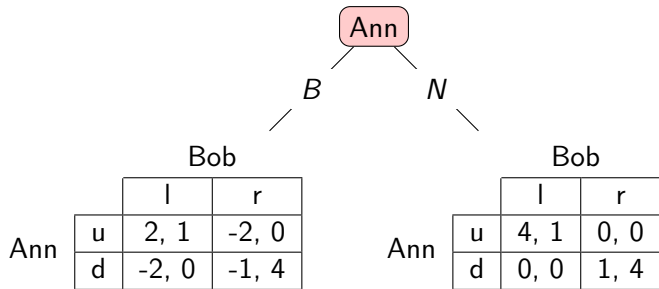
		<i>B</i>	
		l	r
<i>A</i>	<i>Out</i>	2, 1	2, 1
	u	4, 1	0, 0
	d	0, 0	1, 4

		<i>B</i>	
		l	r
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	u	4, 1	0, 0
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		<i>B</i>	
		l	r
<i>A</i>	Out	2, 1	2, 1
	u	4, 1	0, 0
	d	0, 0	1, 4



		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

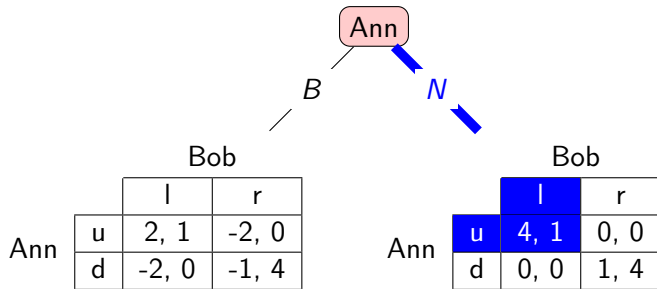
		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
	Nu	4, 1	0, 0	4, 1	0, 0
	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
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	Nd	0, 0	1, 4	0, 0	1, 4

		Bob			
		ll	lr	rl	rr
Ann	Bu	2, 1	2, 1	-2, 0	-2, 0
	Bd	-2, 0	-2, 0	-1, 4	-1, 4
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	Nd	0, 0	1, 4	0, 0	1, 4



What is forward induction reasoning?

Forward Induction Principle: a player should use all information she acquired about her opponents' past behavior in order to improve her prediction of their future simultaneous and past (unobserved) behavior, relying on the assumption that they are rational.

P. Battigalli. *On Rationalizability in Extensive Games*. Journal of Economic Theory, 74, pgs. 40 - 61, 1997.

Thank you!